UNIVERSITY COLLEGE LONDON

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EXAMINATION FOR INTERNAL STUDENTS

For the following qualifications :-

B.Sc. M.Sci.

Mathematics C353: Theory Of Traffic Flow I

COURSE CODE	:	MATHC353
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DATE	:	07-MAY-02
TIME	:	14.30
TIME ALLOWED	:	2 hours

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TURN OVER

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All questions may be answered, but only marks obtained on the best four questions will count. The use of an electronic calculator is not permitted in this examination.

With a conventional notation, the average delay d to a vehicle in a stream of traffic controlled by a fixed-time signal is approximated by

$$d = 0.9\{c(1 - \Lambda)^2/2(1 - y) + x^2/2q(1 - x)\}.$$
(*)

State what is denoted by each quantity that appears in this expression and write down an inequality that is a necessary condition for the approximation to be valid.

At a signal-controlled road junction, there are n traffic streams and quantities relating to stream j are denoted by the relevant symbols from equation (*) with subscript j. The total lost time in the signal cycle is L, $\lambda_0 = L/c$ and equation (*) is valid for every stream.

Show that the average delay per unit time at the junction is approximated by 0.9D, where

$$D = \sum_j \{ f_j(\Lambda_j) / \lambda_0 + g_j(\Lambda_j) \},\$$

where $f_i(\Lambda) = Lq_i(1 - \Lambda)^2/2(1 - y_i)$ and $g_i(\Lambda) = y_i^2/2\Lambda(\Lambda - y_i)$.

Regarding D as a function of the n + 1 variables λ_0 , Λ_1 , Λ_2 , ..., Λ_n , express in terms of λ_0 , the f_j and g_j and their derivatives w.r.t. Λ (which you may write simply as f_j', g_j', f_j'' and g_j'')

(a) $\partial^2 D / \partial \lambda_0^2$ (b) $\partial^2 D / \partial \lambda_0 \partial \Lambda_j$ (c) $\partial^2 D / \partial \Lambda_j \partial \Lambda_k$

Outline briefly without detailed algebra how your results could be used to show that for all suitable signal timings the $(n + 1) \times (n + 1)$ matrix of these second derivatives is positive definite.

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2(a) A signal-controlled road junction has two traffic streams, one having green in stage 1 and the other in stage 2 of a 2-stage cycle. The proportions of the cycle that are effectively green for stages 1 and 2 are λ_1 and λ_2 respectively.

The maximum cycle time is 80 seconds and the total lost time per cycle and the minimum green times for the two stages are each 8 seconds.

Draw in the (λ_1, λ_2) plane the triangle formed by all the points representing signal timings that satisfy the above conditions.

Draw in the same diagram the line corresponding to a cycle time of 40 seconds and let the segment of this line lying in the triangle be AB, where A corresponds to the larger of the values of λ_2 at these two points.

The amounts of traffic in the two streams are such that their degrees of saturation are unity when $\lambda_1 = 0.15$ and $\lambda_2 = 0.30$ respectively.

Without calculating any values of delay, sketch a graph of average delay per unit time at the junction as a function of distance of the point (λ_1, λ_2) along the line *AB* in each direction from an origin at *A*. You may assume that delay is minimised when $\lambda_1 \approx 0.3$.

(b) Customers arrive at a queue according to a Poisson process with mean rate q, and have independent and identical (constant) service times μ . Show that if $x = \mu q < 1$, then the mean delay d incurred by customers in this queue is given by

$$d = \frac{\mu}{2} \left(\frac{x}{1-x} \right).$$

Explain why as it stands, this expression is unsuitable for use to estimate delays at a traffic signal, and describe how it can be used within one that is more suitable.

Wardrop's first principle of route choice is that "the journey times on all routes actually used are equal, and less than those which would be experienced by a single vehicle on any unused route".

In a network in which traffic respects first-in first-out (FIFO) discipline, the demand for travel varies over time. Show that, in order for traffic flows in this network to satisfy Wardrop's first principle, the proportion $\mu_p(s)$ of demand that is assigned to route p at time s satisfies

$$\mu_{p}(s) = \frac{g_{p}[\tau(s)]}{\sum_{q \in P} g_{q}[\tau(s)]} \quad \forall p \in P$$

where $\tau(s)$ is the time of arrival of traffic that departs at time s,

 $g_p(t)$ is the outflow from route p at time t, and

P is the set of routes available for that journey.

Discuss the use in the right-hand side of this expression of route outflows at time $\tau(s)$ to calculate assignment proportions at time $s < \tau(s)$.

4 Explain clearly the physical interpretation and derive from first principles a formula for the speed of each of a *shock wave* and a *traffic wave*. Show that at each density, vehicles travel faster than do waves and establish an expression for the rate at which traffic passes a wave of density k.

A traffic signal is located at position x = 0 on a road, where x increases in the direction of travel and traffic generally moves according to the speed-density relationship

 $v = v_0 \left(1 - \frac{k}{k_j} \right) \quad \left(0 \le k \le k_j \right)$

where v_0 is the free-flow speed and k_j is the jam density. Before time t = -r the signal is green and the traffic flow is uniform everywhere at rate q_a . At that time, the traffic signal changes to red so that a queue of density $k = k_j$ starts to form in the region x < 0. At time t = 0, the signal changes to green and traffic in the queue starts to move according to the speed-density relationship. Show that the maximum upstream extent of the queue occurs when the density k_f at the back of the queue is given by

$$k_{f} = \frac{k_{j}}{2} \left(1 + \sqrt{1 - \frac{4 q_{a}}{v_{0} k_{j}}} \right).$$

Hence or otherwise show that the position x_m and time of occurrence t_m of the maximum upstream position of the queue are given by

$$(x_m, t_m) = \left(-v_0 \sqrt{1 - \frac{4 q_a}{v_0 k_j}}, 1\right) \frac{q_a r}{(v_0 k_j - 4 q_a)}.$$

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Given an $I \times J$ matrix $\tau = (\tau_{ij})$ whose elements are strictly positive and sets of desired strictly positive row and column sums a_i (i = 1, 2, ..., I) and b_j (j = 1, 2, ..., J), define the two kinds of iteration in the Furness iterative procedure whose limit is the unique matrix $\mathbf{t}^* = (t_{ij}^*)$ with the desired row and column sums and such that $t_{ij}^* = r_i s_j \tau_{ij}$

for some r_i and s_j and all i and j.

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A second matrix $\mathbf{\sigma} = (\sigma_{ij})$ has limit \mathbf{s}^* under the Furness procedure. Show that $\mathbf{s}^* = \mathbf{t}^*$ if and only if for some ξ_i and η_j , $\sigma_{ij} = \xi_i \eta_j \tau_{ij}$ for all *i* and *j*.

Suppose that, for all *i* and *j*, $\tau_{ij} = \exp(-\alpha c_{ij})$, where α is a parameter and c_{ij} is the cost of travel from origin zone *i* to destination zone *j* in a city, and that t_{ij}^* is an estimate of the number of journeys from *i* to *j* in a certain period. In a city in which I = 2 and J = 3 and α , the a_i and the b_j are given, which of the following four matrices (c_{ij}) would lead to a different t* from the other three and why?

$\begin{pmatrix} 3\\6 \end{pmatrix}$	4 7	$\begin{pmatrix} 5\\8 \end{pmatrix}$	$\begin{pmatrix} 4\\6 \end{pmatrix}$	5 7	$\begin{pmatrix} 6 \\ 8 \end{pmatrix}$
$\begin{pmatrix} 4\\ 6 \end{pmatrix}$	6 7	$\begin{pmatrix} 6\\ 9 \end{pmatrix}$	$\begin{pmatrix} 5\\7 \end{pmatrix}$	6 8	6 8

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