## Section A. Answer at least ONE question

1. a. i. A binary search tree with $n$ nodes can be represented by four $n$-size arrays. What are these four arrays?
ii. Give a definition for the binary search tree property.
iii. Give a definition for the heap property.
iv. Why is a binary search tree (with the binary search tree property holding) a better data structure for random access than a heap (with the heap property holding)?
b. Draw the heap that corresponds to the following array

$$
\langle 4,1,3,2,16,9,10,14,8,7\rangle
$$

Explain how a sequence of applications of heapify on this heap can give a heap with the heap property holding. Also, give the array corresponding to this new heap.
[8 marks]
c. i. What is a trie and why is it useful for text coding?
ii. Draw the highest trie that can be formed with four leaves, and then give a string of 12 characters that when Huffman's algorithm is applied gives this trie.
iii. Draw the lowest trie that can be formed with four leaves, and then give a string of 12 characters that when Huffman's algorithm is applied gives this trie.
[9 marks]
d. Explain how Huffman's algorithm can be generalized to ternary codewords (ie.code words using the symbols 0,1 , and 2). Also explain why it gives an optimal coding of the input text.
2. See 2b12.q2.ps
3. See 2b12.q3.ps

## Section B. Answer at least ONE question

4. a. Let $L$ be a first-order language, $M=(D, I)$ be an $L$-structure and let $\phi$ be a formula of $L$.
i. Say what it means for $\phi$ to be valid in $M$.
ii. Say what it means for $\phi$ to be satisfiable in $M$.
[8 marks]
b. Let $\phi$ be any formula and let $M$ be a structure. Is the following statement true?

$$
\neg \phi \text { is satisfiable in } M \Longleftrightarrow M \not \vDash \phi
$$

If you think the statement is true then explain why; if not then give a counter-example.
c. Let $L$ be a first-order language with constant symbol $C=\{z\}$, function symbols $F=\{ \}$ and predicate symbols $P=\left\{e q^{2}, g^{2}\right\}$. Let $Z$ be the structure $(\mathbb{Z}, I)$ where $\mathbb{Z}$ is the set of integers, $I=\left(I_{c}, I_{f}, I_{p}\right)$ and

$$
\begin{aligned}
I_{c}(z) & =0 \\
I_{p}\left(e q^{2}\right) & =\{(n, n): n \in \mathbb{Z}\} \\
I_{p}\left(g^{2}\right) & =\{(m, n): m>n\}
\end{aligned}
$$

For each formula below say if the formula is valid in $Z$ and if it is satisfiable in $Z$. If the formula is not valid in $Z$, give an assignment which falsifies the formula. If the formula is satisfiable in $Z$, give an assignment that satisfies it.
i. $\neg\left(\left(g^{2}(x, y) \vee g^{2}(y, x)\right) \vee e q^{2}(x, y)\right)$
ii. $\forall x \forall y\left(\left(g^{2}(x, y) \vee g^{2}(y, x)\right) \vee e q^{2}(x, y)\right)$
iii. $\left(g^{2}(x, y) \rightarrow \exists z\left(g^{2}(x, z) \wedge g^{2}(z, y)\right)\right)$.
iv. $\left(g^{2}(x, z) \wedge \forall y\left(g^{2}(y, z) \rightarrow\left(g^{2}(y, x) \vee e q^{2}(y, x)\right)\right)\right)$.
5. a. What does it mean when we say that a first-order formula is valid?
b. Explain how to use a tableau to discover if a first-order formula is valid or not.
c. For each of the following formulas, construct an appropriate tableau and use this to determine whether the formula is valid.
i. $\exists x \forall y C^{2}(x, y) \rightarrow \forall y \exists x C^{2}(x, y)$
ii. $\exists x\left(P^{1}(x) \vee Q^{1}(x)\right) \rightarrow\left(\exists x P^{1}(x) \vee \exists x Q^{1}(x)\right)$
iii. $\neg[(\exists x \exists y(x<y)) \wedge \forall x \forall y(x<y \rightarrow \exists z(x<z \wedge z<y))]$

