Where an algorithm is asked for, you may write in any suitable pseudocode. Correct syntax for any computer language is not expected.

## Answer 3 questions.

 You have been asked to visit a local school to give the sixth-form computer science students a two-hour lecture on distributed systems. This must include an outline of whether and how you believe distributed systems to be important, what problems the designers of distributed systems must deal with, and whether such problems are fundamental or can be overcome.

Write a detailed outline for this lecture, saying what aspects of distributed systems you have chosen to explore, why, and how each sub-topic relates to those around it and to the overall goals you have for the lecture.

[33]

2)

- a) Digital signatures have different properties to traditional signatures placed on a document with a pen. What is a digital signature and what are the functional differences between them and those written in ink?
- b) In PGP, an MD5 hash value is used to produce digital signatures. Why are hash values used, and what properties do we require of such a hash value. [5]
- c) What threats are there to data being passed over insecure links, and how can they be avoided? [12]
- d) Present and explain an algorithm by which a user may authenticate themselves to a computer system using public key cryptography. You should show that your algorithm avoids the threats you identified in c), or else demonstrate that this is not possible.

[12]

[TURN OVER]

- 3)
- a) Explain the operation of a multilevel paging system. Use as an example a system with a 32 bit address, 4K word pages, and a 2 level paging system with page tables of equal size at each level. Your answer should include details of the sizes of page tables and the way that you calculated these
- b) Given the above system and assuming a memory size of 8 Mbyte, explain what happens when the following program is run.

```
int array[4096][4096]
for (int i = 0; i < 4096; i++)
for (int j = 0; j < 4096; j++)
array[i][j]++</pre>
```

c) What happens when, instead, the following program is run:

```
int array[4096][4096]
for (int j = 0; j < 4096; j++)
for (int i = 0; i < 4096; i++)
array[i][j]++ [3]</pre>
```

- d) If all pages are initially on disk, memory access time is 200ns, and disk access time
   = 2ms + (250ns/byte transferred), what is the time to run each of the programs
   above programs? State any assumptions you make
- e) What is a TLB and how could one improve the running time in the above cases? Assume a TLB access time of 1ns and a 512 Kbyte cache size, what running times would one now expect for the above programs? Comment on your results.

[CONTINUED]

[3]

[11]

[10]

[6]

- a) What three criteria must be satisfied in any solution to the critical section problem, and why?
- b) The following are two attempts at solving the two process critical section problem, using only an assumption that read and assignment operations are atomic. With reference to your answer to part a), say how the solutions are deficient.

```
//
// Solution A
11
int turn = 1
               // Global variable shared by both processes
process pl:
                                                      process p2:
while (TRUE)
                                                      while (TRUE)
{ while (turn == 2)
                                                      { while (turn == 1)
    no_op
                                                          no_op
  CRITICAL SECTION
                                                        CRITICAL SECTION
  turn = 2
                                                        turn = 1
  NON CRITICAL SECTION
                                                        NON CRITICAL SECTION
}
                                                      }
11
// Solution B
11
int c1 = FALSE
                 // Global variables shared by both processes
int c2 = FALSE
process pl:
                                                     process p2:
while (TRUE)
                                                      while (TRUE)
                                                      \{ c2 = TRUE \}
\{ c1 = TRUE \}
                                                        while (c1 == TRUE)
{ c2 = FALSE
  while (c2 == TRUE)
  { c1 = FALSE
    sleep (10 milliseconds)
c1 = TRUE
                                                          sleep (10 ms)
c2 = TRUE
  }
                                                        }
  CRITICAL SECTION
                                                        CRITICAL SECTION
  c1 = FALSE
                                                        c2 = FALSE
  NON CRITICAL SECTION
                                                        NON CRITICAL SECTION
}
                                                      }
```

c) Using only the assumption that read and write operations are atomic, give a correct solution to the 2-process critical section problem and say why it is correct.

[11]

[6]

[16]

[TURN OVER]

4)

a) What metrics could one use to measure the performance of a scheduling algorithm?

[5]

[16]

- . ....
- i) What is the convoy effect and why might it be a problem? Illustrate your answer with an example. [4]
- ii) Show how a different scheduling algorithm of your choice overcomes the problem. [4]
- iii) Is this second scheduling algorithm better in all circumstances? Justify your answer. [4]
- c) On a particular computer, the operating system has a context switch time of 1ms, and a preemptive scheduling system with three priority levels, A, B, and C. Each priority level has a separate queue associated with it, all scheduled using round robin, with the constraint that if there is a process in the level A queue it is scheduled in preference to anything in levels B and C, and a process in level B is scheduled in preference to anything in level C. In level A, the timeslice is 20ms, in level B it is 40ms, and in level C it is 80ms. Processes initially arrive at level A but their priority may change according to the following two rules:
  - i) If a process uses its entire timeslice, it sinks one level, unless already at level C when it remains there.
  - ii) After 100ms, any process which has not run in the previous 100ms is promoted by one level, unless already in level A.
  - iii) After returning from i/o, a process re-enters the level A queue.
  - iv) If a process is preempted, it does not change its place in the queue.

Given the process arrivals below, and assuming that each process executes in an endless sequence of CPU and i/o bursts, draw a Gantt chart to show how this system performs over the first 202ms.

Process	<b>P</b> <sub>0</sub>	<b>P</b> <sub>1</sub>	<b>P</b> <sub>2</sub>	<b>P</b> <sub>3</sub>	<b>P</b> <sub>4</sub>
CPU burst	5	10	50	100	5
i/o burst	20	10	10	10	50
Arrival time	0	10	20	30	40

[END OF PAPER]

b)