

**Section A.**

1.  $a \in R$  is a *unit* if there exists  $b \in R$  with  $ab = ba = 1$ .

[2 marks]

$b$  is then the *multiplicative inverse* of  $a$ . (It is unique.)

[2 marks]

a)  $\mathbf{Z}_3 = \{0, 1, 2\}$ .  $1 \times 1 = 2 \times 2 = 1$  (and  $0 \times a = 0$  for all  $a \in \mathbf{Z}_3$ ). So 1 and 2 are the units in  $\mathbf{Z}_3$ .

[2 marks.]

b)  $\mathbf{Z}_4 = \{0, 1, 2, 3\}$   $1 \times 1 = 3 \times 3 = 1$  and  $2 \times 2 = 0$ . No element can be both a unit and a zero divisor. So 1 and 3 are the units in  $\mathbf{Z}_4$ .

c) 
$$(a_n x^n + \dots)(b_m x^m + \dots) = a_n b_m x^{n+m} + \dots$$

So the only way in which to get product 1 is if both polynomials are constant. So the units are all the constant polynomials  $a_0$  with  $a_0 \neq 0$ , with inverse  $1/a_0$ .

[3 marks.]

[10 = 2 + 1 + 2 + 2 + 3 marks.]

The first part is standard theory, the rest standard homework problems.

2.  $a(x) = x^4 + x^2 + 1$ ,  $b(x) = x^3 + 1 \in \mathbf{Z}_2[x]$ . We have

$$x^4 + x^2 + 1 = x(x^3 + 1) + x^2 + x + 1 \tag{1}$$

by inspection or long division

[2 marks] and

$$x^3 + 1 = (x + 1)(x^2 + x + 1) \tag{2}$$

again by inspection or long division

[2 marks]

So  $c(x) = x^2 + x + 1$  is the g.c.d. of  $a(x)$  and  $b(x)$

[1 mark] and  $d(x) = x + 1$  from (2), while from (1) and (2) (or alternatively by long division)

$$x^4 + x^2 + 1 = x(x + 1)(x^2 + x + 1) + x^2 + x + 1 = (x^2 + x + 1)(x^2 + x + 1)$$

and  $e(x) = x^2 + x + 1$ .

[2 marks]

From (1),

$$x^2 + x + 1 = x^4 + x^2 + 1 + x(x^3 + 1).$$

So  $m(x) = 1$  and  $n(x) = x$ .

[1 mark]

[8 = 2 + 2 + 1 + 2 + 1 marks]

Standard homework exercise.

3.  $x^4 + 1$  has no zeros in  $\mathbf{Z}$  (or  $\mathbf{R}$ ).

[1 mark.]

So if it factorizes in  $\mathbf{R}[x]$  or  $\mathbf{Z}[x]$ , it must have a factorization as a product of two degree two polynomials, and we are allowed to assume this is of the form

$$(x^2 + ax + b)(x^2 - ax + b).$$

(This actually follows from considering the coefficients of the  $x^3$  term and then the  $x$  term, since we must have  $a \neq 0$ .) Then  $b^2 = 1$ , giving  $b = \pm 1$ . Then  $a^2 = 2b = \pm 2$ , giving no real  $a$  unless  $b = 1$ , in which case  $a = \sqrt{2}$ . Both factors have no real zeros because  $(\sqrt{2})^2 - 4 < 0$ , and hence they are irreducible in  $\mathbf{R}[x]$ .

[5 marks]

So there is a factorization in  $\mathbf{R}[x]$

$$x^4 + 1 = (x^2 + \sqrt{2}x + 1)(x^2 - \sqrt{2}x + 1)$$

but  $x^4 + 1$  is irreducible in  $\mathbf{Z}[x]$  because  $\sqrt{2} \notin \mathbf{Z}$ .

[2 marks]

In  $\mathbf{Z}_2[x]$  we have

$$x^4 + 1 = (x + 1)^4,$$

which is the factorization into irreducibles in  $\mathbf{Z}_2[x]$ .

[3 marks]

[11 = 1 + 5 + 2 + 3 marks]

Standard homework exercise.

4a)  $x + y + 2 = 0 \Leftrightarrow y = 1 + 2x$  (in  $\mathbf{Z}_3$ ).

So taking  $x = 0, 1, 2$  gives the points  $(0, 1), (1, 0), (2, 2)$ .

b)

$$x + y + 2 = 0, 2x + y + 1 = 0.$$

Subtracting the first equation from the second gives  $x - 1 = 0$ , that is,  $x = 1$ , so that from the first equation  $y = 1 + 2x = 0$ . So  $(1, 0)$  is the point of intersection.

[2 marks.]

c) The lines through  $(1, 0)$  are of the form

$$a(x - 1) + by = 0$$

for  $(a, b) \neq (0, 0)$ . To avoid duplication of lines, we take  $b = 1$  or  $b = 0$  and  $a = 1$ .

$b = 1$  gives  $y = 0, x + y + 2 = 0, 2x + y + 1 = 0$ .

$b = 0$  and  $a = 1$  gives  $x + 2 = 0$ .

[3 marks]

[7 = 2 + 2 + 3 marks.]

[Standard homework exercise.]

5a) 6 does not divide  $11 \times 3$ . So there is no 1-design with these parameters.

[1 mark]

b)  $6 \mid 10 \times 3$  and  $3 < \binom{9}{5}$ . So there is a 1-design with these parameters.

[2 marks]

c) The set of lines in  $P^2(\mathbf{Z}_2)$  is such a 2-design:  $P^2(\mathbf{Z}_2)$  has 7 points, each line has 3 points and any two lines intersect in 1 point.

[3 marks]

d) There is no 2 design with these parameters  $(v, k, r)$  because  $2 = k - 1$  does not divide  $(v - 1)r = 7$ .

[3 marks]

[9 = 1 + 2 + 3 + 3 marks]

Standard homework exercises.

6. All the columns of  $H$  are distinct and not identically 0. So  $H$  corrects one error

[2 marks]

The weight has to be 3

[1 mark.]

a)

$$\begin{pmatrix} 0 & 1 & 0 & 1 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 \end{pmatrix} \begin{pmatrix} 1 \\ 1 \\ 1 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 1 \end{pmatrix}.$$

This is not a code word. Since the third column of the check matrix is obtained, the third letter must be wrong and the corrected word is 1100110.

[4 marks]

b)

$$\begin{pmatrix} 0 & 1 & 0 & 1 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \\ 0 \\ 1 \\ 0 \\ 1 \\ 0 \\ 1 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}.$$

So this is a codeword

[3 marks.]

10 = 2 + 1 + 4 + 3 marks.]

Standard homework exercise, with a bit of standard theory at the start.

### Section B.

7(i)  $S \subset R$  is a *subring* of  $R$  if  $S \neq \phi$  and  $x, y \in S \Rightarrow x - y \in S, xy \in S$ .

[2 marks]

$S \subset R$  is an *ideal* of  $R$  if  $S \neq \phi$  and  $x, y \in S \Rightarrow x - y \in S, x \in S, y \in R \Rightarrow xy \in S$ .

(i)a)  $(2 + i) = -1 + 2i \notin S$ , although  $2 + i, i \in S$ . So  $S$  is not a subring, and hence not an ideal

[2 marks]

$$(i)b) \quad (m_1 + 2n_1i) - (m_2 + 2n_2i) = (m_1 - m_2) + 2(n_1 - n_2)i.$$
$$(m_1 + 2n_1i)(m_2 + 2n_2i) = (m_1m_2 - 4n_1n_2) + 2(n_1m_2 + n_2m_1)i$$

for all integers  $m_i, n_i, i = 1, 2$ . So  $S$  is a subring.

[3 marks]

But  $S$  is not an ideal because  $1 \in J$  but  $i \cdot 1 = i \notin J$ .

[1 mark]

(ii) Let  $b \neq 0, b \in J$  with  $v(b)$  minimal. Take any  $a \in J$ . Then  $a = qb + r$  for some  $a, r \in R$  with  $r = 0$  or  $v(r) < v(b)$ . We have  $r = a - qb \in J$ . Since  $v(b)$  is minimal for  $b \neq 0$ , we must have  $r = 0$ , and hence  $a = qb$ . So  $J \subset bR$ . But  $bR \subset J$ . So  $J = bR$ .

[5 marks.]

[15 = 2 + 2 + 2 + 3 + 1 + 5 marks]

(i) is standard theory for the definitions and similar to homework exercises for a) and b), while (ii) is standard theory, with hints provided.

8(i). For  $f(x) = x^3 + 2x + 1$ , we have  $f(0) = 1, f(1) = 1, f(2) = 13 = 1$ . So  $f$  has no zeros in  $\mathbf{Z}_3$ , hence no degree one factors, and is irreducible

[2 marks]

$$f(x+1) = (x+1)^3 + 2(x+1) + 1 = x^3 + 2x + (3x^2 + 3x) + 1 + 2 + 1 = f(x).$$

[2 marks]

$$f(x+2) = (x+2)^3 + 2(x+2) + 1 = x^3 + 6x^2 + 12x + 8 + 2x + 4 + 1 = f(x).$$

[1 mark]

(ii)

$$f(J+x) = J + x^3 + 2(J+x) + 1 = J + x^3 + 2x + 1 = J = 0.$$

[1 mark]

$a$  is a zero  $\Leftrightarrow a+1$  and  $a+2$  are zeros by a) (since  $a = (a+1) + 2 = (a+2) + 1$ ). So  $J+x+1$  and  $J+x+2$  are the other zeros.

[1 mark]

(iii)

$$J+g_1(x) = J+g_2(x) \Leftrightarrow f(x) \mid (g_1(x)-g_2(x)) \Leftrightarrow f(x) = f(x+1) \mid (g_1(x+1)-g_2(x+1))$$
$$\Leftrightarrow J+g_1(x+1) = J+g_2(x+1).$$

[3 marks]

This shows that  $\varphi(J+g(x)) = J+g(x+1)$  is well-defined and one-to-one.

[1 mark]

To show that it is onto and compute the inverse:  $\psi(J+h(x)) = J+h(x+2)$  is also well-defined for similar reasons (or alternatively,  $\psi = \varphi \circ \varphi$ ) and

$$\psi \circ \varphi(J+g(x)) = J+g(x+3) = J+g(x) = \varphi \circ \psi(J+g(x)).$$

[2 marks.]

(iv) Any isomorphism  $\Phi$  has to map  $J+x$  to a zero of  $f$ , that is, to  $J+x$  or  $J+x+1$  or  $J+x+2$ , and then  $\Phi(J+g(x)) = g(\Phi(J+x))$ .

[2 marks]

[15 = 2 + 2 + 1 + 1 + 1 + 3 + 1 + 2 + 2 marks.]

Standard exercises on quotient rings, parts of (iii) and (iv) are basically unseen, but also similar in part to standard homework exercises on computing inverse homomorphisms.

9a) Use the 2-design of lines in  $\mathbf{Z}_3^2$ . The 9 points in  $\mathbf{Z}_3^2$  correspond to the cars. The sessions correspond to the sets of parallel lines. Each set of parallel lines is of the form  $\{ax + by + c = 0 : c \in \mathbf{Z}_3\}$  for fixed  $(a, b) \neq (0, 0)$  (where  $(\lambda a, \lambda b)$  gives the same line as  $(a, b)$  for  $\lambda \neq 0$ ). The test-drivers correspond to the values of  $c \in \mathbf{Z}_3$ . In the incidence matrix below, 1 indicated that a car is tested by that driver in that session.

[4 marks for explanation]

*	(0, 0)	(1, 0)	(2, 0)	(0, 1)	(1, 1)	(2, 1)	(0, 2)	(1, 2)	(2, 2)
Session 1	*	*	*	*	*	*	*	*	*
$x = 0$	1	0	0	1	0	0	1	0	0
$x + 1 = 0$	0	0	1	0	0	1	0	0	1
$x + 2 = 0$	1	0	0	1	0	0	1	0	
Session 2	*	*	*	*	*	*	*	*	*
$y = 0$	1	1	1	0	0	0	0	0	0
$y + 1 = 0$	0	0	0	0	0	0	1	1	1
$y + 2 = 0$	0	0	0	1	1	1	0	0	0
Session 3	*	*	*	*	*	*	*	*	*
$x + y = 0$	1	0	0	0	0	1	0	1	0
$x + y + 1 = 0$	0	0	1	0	1	0	1	0	0
$x + y + 2 = 0$	0	1	0	1	0	0	0	1	
Session 4	*	*	*	*	*	*	*	*	*
$x + 2y = 0$	1	0	0	0	1	0	0	0	1
$x + 2y + 1 = 0$	0	0	1	1	0	0	0	1	0
$x + 2y + 2 = 0$	0	1	0	0	0	1	1	0	0

[7 marks]

b) The 2-design corresponding to Kirkman's Schoolgirls problem will do. The 7 days in Kirkman's problem correspond to the 7 testdriving session. The 15 girls correspond to the 15 cars. The 5 groups of girls (on each day) correspond to the 5 groups of cars driven by the 5 testdrivers.

4 marks]

[15 = 4 + 7 + 4 marks]

Part a) is similar to standard homework exercise, b) tests knowledge of Kirkman's schoolgirls problem (and would be similar to a standard homework problem, but harder, if it were developed).

10(i) We have  $1 + x^3 = (1 + x)(1 + x + x^2)$  in  $\mathbf{Z}_2[x]$ . Hence, since  $(1 + y)^2 = 1 + y^2$ , we have

$$1 + x^{12} = (1 + x^3)^4 = (1 + x)^4(1 + x + x^2)^4.$$

$1 + x + x^2$  is irreducible in  $\mathbf{Z}_2[x]$  since it has no zero in  $\mathbf{Z}_2$  - and of course  $1 + x$  is irreducible as well

[3 marks]

(ii) A general factor of  $1 + x^{12}$  is of the form  $(1 + x)^r(1 + x + x^2)^s$  where  $0 \leq r, s \leq 4$ . To get a factor  $g(x)$  of degree 5 we need  $r + 2s = 5$ , and the only possibilities are  $r = 3, s = 1$  or  $r = 1, s = 2$ . To get a factor  $h(x)$  of degree 7 we must have  $r = 3, s = 2$  or  $r = 1, s = 3$ .

[2 marks]

a) With canonical generator  $g(x) = (1 + x)^3(1 + x + x^2)$ , the check matrix is computed from

$$h(x) = (1 + x)(1 + x + x^2)^3 = (1 + x^3)(1 + x^2 + x^4) = x^7 + x^5 + x^4 + x^3 + x^2 + 1$$

from which the check matrix is

$$\begin{pmatrix} 1 & 0 & 1 & 1 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 1 \end{pmatrix}.$$

All columns are distinct and nonzero. So this code corrects one error.

[5 marks.]

b) With canonical generator  $g(x) = (1 + x)(1 + x + x^2)^2$ , the check matrix is computed from

$$h(x) = (1 + x)^3(1 + x + x^2)^2 = (1 + x)(1 + x^6) = x^7 + x^6 + x + 1,$$

from which the check matrix is

$$\begin{pmatrix} 1 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 1 \end{pmatrix}$$

The first and seventh columns are the same. So the code has weight two and does not correct one error.

5 marks

[15 = 3 + 2 + 5 + 5 marks.]

Standard homework exercise.

11(i) Every row of  $A$  has exactly three 1's (because there are three points on every line in  $P^2(\mathbf{Z}_2)$ ). For any two rows of  $A$ , there is exactly one column where both have a 1, because any two lines intersect in exactly one point. So the distance between any two rows of  $A$  is 4.

[4 marks.]

Write  $\mathbf{1}, \mathbf{0}$  for the vectors of all 1's and all 0's. If  $\mathbf{c}$  and  $\mathbf{c}'$  are rows of  $A$  with  $\mathbf{c} \neq \mathbf{c}'$ , then  $\mathbf{c}$  and  $\mathbf{1} - \mathbf{c}'$  agree in the entries where  $\mathbf{c}$  and  $\mathbf{c}'$  differ. They agree in exactly four entries, and hence differ in three.

[3 marks.]

Let  $E$  be the matrix of all 1's. Since any row of  $A$  or  $E - A$  has either 3 or 4 1's, these all differ from  $\mathbf{0}$  and  $\mathbf{1}$  in at least three entries. So the minimum distance of  $C$  is 3.

[2 marks.]

(ii) There are  $\binom{v}{i}$  words which differ from  $\mathbf{c}$  in exactly  $i$  entries (this being the number of  $i$ -element subsets of a  $v$ -element set). So the number of words differing from  $\mathbf{c}$  in  $\leq e$  entries is

$$1 + \binom{v}{1} + \cdots + \binom{v}{e}.$$

[3 marks.]

Write  $B(\mathbf{c})$  for this set. For minimum distance  $2e + 1$ , the sets  $B(\mathbf{c})$  must all be disjoint. Since  $C'$  has  $m$  words and  $\mathbf{Z}_2^v$  has  $2^v$  words, this gives

$$2^v \geq m \left( 1 + v + \cdots + \binom{v}{e} \right).$$

[2 marks.]

To get equality, take  $C'$  to be the code  $C$  of part a), which has  $m = 16$  words in  $\mathbf{Z}_2^7$  and has minimum distance  $2e + 1 = 3$  with  $e = 1$ . The equality then says  $2^7 = 16 \times (7 + 1) = 2^4 \times 2^3$  - which is true.

1 mark.

[15 = 4 + 3 + 2 + 3 + 2 + 1 marks.]

(i) is similar to a homework exercise. Most of (ii) is basic theory. The last bit is unseen.