

Solutions for May exam

1. A zero-divisor in a ring R is an element $r \in R$ such that $r \neq 0$ and there is an $s \neq 0 \in R$ such that $rs = 0$. [lecture] (1 mark)
 - (a) 6 is a zero-divisor in $\mathbf{Z}/8$ because $6 \cdot 4 = 0$ in $\mathbf{Z}[8]$. [lecture] (2 marks)
 - (b) 4 is not a zero-divisor in $\mathbf{Z}/11$ because \mathbf{Z}/p has no zero-divisors when p is prime. [lecture] (2 marks)
 - (c) $3x^2 + 5x + 7/3$ is not a zero-divisor in $\mathbf{Q}[x]$ because \mathbf{Q} has no zero-divisors, and we know that if R has no zero-divisors neither does $R[x]$. [similar to lecture] (2 mark)
 - (d) $x - 1$ is a zero-divisor in $\mathbf{Z}/2[x]/(x^5 - 1)$ because $(x - 1)(x^4 + x^3 + x^2 + x + 1) = 0$ in $\mathbf{Z}/2[x]/(x^5 - 1)$. (It is acceptable to explain this by saying that $x - 1$ divides $x^5 - 1$.) [similar to lecture] (2 marks)
- 2a. We have $d(r) = 45, d(s) = 50$, so let's divide s by r . This gives $s/r = (7+i)/(-3+6i) = (-15-45i)/45 = -1/3 - i$. The closest element of $\mathbf{Z}[i]$ to this is $-i$, so we have $7 + i = (-i)(-3 + 6i) + 1 - 2i$ and $\gcd(r, s) = \gcd(1 - 2i, -3 + 6i)$. Since $-3 + 6i = -3(1 - 2i)$, the GCD is $1 - 2i$. Also from the previous equation we have $7 + i + i(-3 + 6i) = 1 - 2i$. [similar to lecture and homework] (4 marks)
- 2b. Choosing to divide s by r , we get $s = 2r + 4x^2 + x + 4$, then $r = (4x + 2)(4x^2 + x + 4) + 3$, and since 3 is a unit it is a GCD. Working backwards to find a and b , we get successively $3 = r + (x + 3)(4x^2 + x + 4)$, then $3 = r + (x + 3)(s - 2r)$, and so $3 = 3xr + (x + 3)s$. [similar to lecture and homework] (5 marks)
- 3a. It is sufficient to remove all linear factors, since $\deg f = 3$. We see that 1 is a root of f , so $x + 2$ divides f , and indeed $f = (x + 2)(2x^2 + 2)$ in $\mathbf{Z}/3[x]$. Letting $g(x) = 2x^2 + 2$, we calculate $g(0) = 2, g(1) = g(2) = 1$, so g has no roots and is therefore irreducible (its degree being less than 4). Thus $(x + 2)(2x^2 + 2)$ is the factorization. [similar to lecture and homework] (3 marks)
- 3b. Again, it's enough to remove all linear factors, but this time one checks that $f(0) = 4, f(1) = 1, f(2) = 1, f(3) = 1, f(4) = 3$. That is, f has no roots mod 5, so it is irreducible in $\mathbf{Z}/5[x]$. [similar to lecture and homework] (3 marks)
- 3c. From (b), f is irreducible in $\mathbf{Q}[x]$, so the only possible factorization is to divide through by an integer, and indeed $2|f$. So the irreducible factorization is $2(x^3 - x^2 + x + 2)$. [similar to lecture and homework] (3 marks)
- 4i. The required table is

·	0	1	i	$i + 1$
0	0	0	0	0
1	0	1	i	$i + 1$
i	0	i	1	$i + 1$
$i + 1$	0	$i + 1$	$i + 1$	0

- (It is acceptable to choose different coset representatives.) [similar to lecture and homework] (5 marks)
- 4ii. A field is a ring with more than one element in which every nonzero element is a unit (or, has a multiplicative inverse, or, in which for every $r \neq 0 \in R$ there is $s \in R$ with $rs = 1$). This ring is not a field because $1 + i$ is not a unit. (It is a sufficient explanation to say that $1 + i$ is a zero-divisor.) [similar to lecture and homework] (4 marks)
 5. A t -(v, k, r)-design is a set of subsets T_i of size k of a set S of size v such that each t -element subset of S is a subset of exactly r of the T_i . [lecture] (3 marks) If $\{T_i\}$ is a 1-design, then each element of S is an element of exactly r of the T_i , so the sum of the cardinalities of the T_i is vr . On the other hand, the T_i are sets of size k and so $k|vr$. Also there are only $\binom{v-1}{k-1}$ subsets of S of size k that contain a given element (it's not necessary to explain why), so r may not be greater than this. [lecture] (6 marks)
 - 6i. This code has dimension $5 - 3 = 2$ and its information rate is therefore $2/5$. [similar to lecture and homework] (4 marks)
 - 6ii. To list the words, we solve the equations

$$\begin{aligned} x_1 + x_4 + x_5 &= 0 \\ x_3 + x_5 &= 0, \\ x_1 + x_2 + x_3 + x_5 &= 0 \end{aligned}$$

finding the general solution x_4, x_5 arbitrary, $x_3 = x_5, x_2 = x_4 + x_5, x_1 = x_4 + x_5$. Thus the codewords are 00000, 11010, 11101, 00111. From this list the minimum distance is 3 (this is the minimum number

of 1's in any nonzero word) and so the code corrects $\lfloor \frac{3-1}{2} \rfloor = 1$ error. [similar to lecture and homework] (6 marks)

7. An ideal is a nonempty subset I of a ring R such if $a, b \in I$ and $r \in R$ then $a + b$ and $ar \in R$. (Minor variations of this definition are possible.) A principal ideal is an ideal consisting of the multiples of a single element of R . [lecture] (3 marks)

- a. This is an ideal. This can be seen either from the fact that it is the kernel of the homomorphism $\mathbf{Z}[x] \mapsto \mathbf{Z}/3$ given by reducing the constant coefficient mod 3, or more directly as follows: let $f = \sum a_i x^i$ and $g = \sum b_j x^j \in I$, so that $3|a_0$ and $3|b_0$. Then the constant coefficient of $f + g$ is $a_0 + b_0$, a multiple of 3, and for any polynomial h with constant coefficient c_0 the constant coefficient of fh is $a_0 c_0$, again a multiple of 3. [homework] (3 marks)
- b. This is not an ideal, because $3 \in I$ and $i \in R$ but $3i \notin R$. [similar to homework] (3 marks)
- c. This is not an ideal, because $4 \in I$ and $9 \in I$ but $4 + 9 \notin I$. [unseen] (3 marks)
- d. This is an ideal. Again this can be seen directly, by taking $a_0 + b_0 i$, $a_1 + b_1 i$ in the ideal and $c + di$ in the ring and noticing that $a_0 + b_0 i + a_1 + b_1 i = (a_0 + a_1) + (b_0 + b_1)i$, so that $(a_0 + a_1) + (b_0 + b_1)i = (a_0 + b_0) + (a_1 + b_1)i$, the sum of two even numbers, is even. Then $(a_0 + b_0 i)(c + di) = (a_0 c - b_0 d) + (a_0 d + b_0 c)i$, and $a_0 c - b_0 d + a_0 d + b_0 c = (a_0 + b_0)(c + d) - 2b_0 d$ is also even. This time it is definitely easier to notice that this set is the kernel of the homomorphism $\phi : \mathbf{Z}[i] \mapsto \mathbf{Z}/2$ defined by $\phi(a + bi) = a + b \pmod{2}$. [similar to lecture and homework] (3 marks)

- 8 A prime ideal is an ideal I of a ring R such that if $a, b \in R$ and $ab \in I$ then $a \in I$ or $b \in I$ and I is not all of R . (3 marks)

If I is prime, then I is not all of R , so R/I is nontrivial. In addition, if I is prime and $(c + I)(d + I) = 0$ for $c + I, d + I \neq 0 \in R/I$, then $c, d \notin I$ but $cd \in I$, contradiction. Together this proves that R/I is an integral domain. (6 marks)

If R/I is an integral domain, it is nontrivial, so I is not all of R . In addition, if $cd \in I$ for $c, d \notin I$, then $(c + I)(d + I) = 0$ in R/I but $c + I, d + I \neq 0$, which contradiction proves that I must be prime. (6 marks) [lecture]

Omitting the two bits on nontriviality only costs marks once.

- 9i. R/I has $3^3 = 27$ elements, and, because R/I is a field, all but one of those is an element of $(R/I)^*$, so there are 26. The possible orders of these elements are the divisors of 26, namely 1, 2, 13, 26. [similar to lecture and homework] (6 marks)
- 9ii. Indeed, $(x^2 + 1)^2 = x^4 + 2x^2 + 1 = xf + x + 1$ in R , so it is $x + 1$ in R/I . [similar to lecture and homework] (3 marks)
- 9iii. Clearly the order is not 1 or 2, and $(x^2 + 1)^{26} = 1$ in R/I . So $(x + 1)^{13} = (x^2 + 1)^{26} = 1$ in R/I , and the order is 13. [similar to homework] (6 marks)
- 10a. The lines on a projective plane over a field of 2 elements give a 2-(7, 3, 1)-design. Abbreviating the elements as 001, 010, 011, 100, 101, 110, 111, the sets are (001, 010, 011), (001, 100, 101), (001, 110, 111), (010, 100, 110), (010, 101, 111), (011, 100, 111), and (011, 101, 110). [lecture] (8 marks)
- 10b. Similarly, the lines an affine plane over a field of 8 elements give a 2-(64, 8, 1)-design. Such a field exists because 8 is a power of 2 (more of an explanation is not necessary). (6 marks) [lecture]
- 5 marks for checking that these designs don't violate the numerical constraints.
- 11i. Indeed, the quotient is $x^3 + x + 1$. [lecture and homework] (2 marks)
- 11ii. The check matrix is formed by writing rows which represent the product of the quotient from (i) by powers of x , written in order of decreasing powers: thus

$$M = \begin{pmatrix} 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 0 & 0 & 0 & 0 \end{pmatrix}.$$

[lecture and homework] (4 marks)

- 11iii. The desired theorem is that no column of the check matrix should be a multiple of any other, which is easily seen to be true in this case. [lecture and homework] (3 marks)

11iv. To decode the word v , we compute Mv^t , getting respectively 01011, 00000, and 00002. Thus the first word is wrong in the place corresponding to the column that is a multiple of 01011, that is, the fourth; the word should be 21011202. The second word is correct, and the third word is wrong in the first column; it should be 20221011. [lecture and homework] (2 marks each)