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THE UNIVERSITY of LIVERPOOL

JANUARY 2007 EXAMINATIONS

Bachelor of Science : Year 2 No qualification aimed for : Year 1

Decision, Computation and Language

TIME ALLOWED : 2 Hours

INSTRUCTIONS TO CANDIDATES

Answer **all** questions in Section 1 and **two** questions from Section 2. Section 1 accounts for 60% of credit and Section 2 accounts for the remaining 40%.

If you attempt to answer more questions than the required number of questions (in any section), the marks awarded for the excess questions will be discarded (starting with your lowest mark).



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Section 1

Answer all questions in this section.

1.

a. For each of the following pairs of regular expressions over alphabet {a, b}, say whether they represent the same language, and if they do not, give an example of a word that belongs to one language but not the other.

i. ${ab, aab}({ab, aab})^*$ and $({ab, aab})^*aab$	[2 marks]
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- ii. $(\{ab, aab\})^*$ and $(\{ab, (aab)^*\})^*$ [2 marks]
- iii. $(\{b, ba\})^* \cup (\{ba, baa\})^*$ and $(\{b, ba, baa\})^*$ [2 marks]
- b. Give an example of a language that is context-free but is not regular. Write down a context-free grammar for the language. [5 marks]
- c. Consider the following regular grammar, with start symbol S and additional variables A and B.

S	\longrightarrow	cS	
S	\longrightarrow	cA	
A	\longrightarrow	cB	
A	\longrightarrow	a	
B	\longrightarrow	Ъ	

Find a regular expression that represents the language generated by the above grammar. (You may find it helpful to simplify the grammar via back substitution of variables.)

[4 marks]

2.

- a. Draw a diagram of a deterministic finite automaton that accepts the language (over alphabet {a, b, c}) given by the regular expression c*(a* ∪ b*). [5 marks]
- b. Explain how your DFA of part (a) can be altered to obtain a new DFA that accepts the complement of the language of the original DFA (that is, the set of all words not accepted by the original one).
- c. Describe a general method for modifying a deterministic finite automaton M so that if L is the language accepted by M, the modified finite automaton (which may be nondeterministic) accepts L^* , the closure of L. [5 marks]



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3. A language L is said to be prefix-closed if, given any word $w \in L$, all prefixes of w are also members of L. For example, any prefix-closed language that contains the word cats must also contain the words cat, ca, c and the empty word ϵ .

Consider the languages $L_1 = \{a^n b^m : n > m\}$, and $L_2 = \{b^m a^n : n > m\}$. (So, the first of these, L_1 , is strings consisting of a's followed by b's, where there are more a's than b's. L_2 is b's followed by a's, where again there are more a's than b's.)

a.	Which of these is prefix-closed and why?	[5 marks]
b.	Use the pumping lemma to prove that L_2 is not regular.	[5 marks]
c.	Is L_1 regular? Give a brief explanation for your answer.	[5 marks]

4.

- a. Explain what is meant by a *recursive* language, and what is meant by a *recursively enumerable* language. [5 marks]
- b. Suppose that two languages L_1 and L_2 are accepted by Turing machines M_1 and M_2 . Explain how M_1 and M_2 can be used to construct a Turing machine M that accepts words that belong to both L_1 and L_2 . [10 marks]

Section 2

Answer two questions in this section.

5.

- a. Give an unambiguous context-free grammar that can generate words of the form a*b*c* (strings consisting of a sequence of a's followed by a sequence of c's). Explain why your grammar is unambiguous. [5 marks]
- b. Recall Chomsky Normal Form, in which all rules of a context-free grammar are of the form S → ϵ, X → YZ, or X → a, where S is the starting symbol, X, Y and Z may be any variable symbol, and a may be any alphabet symbol. Convert the following grammar to Chomsky Normal Form. You may use the general algorithm, or other transformations.

$$\begin{array}{rccc} S & \longrightarrow & \mathrm{abc}A \\ A & \longrightarrow & \mathrm{a}B\mathrm{a}B \\ B & \longrightarrow & C \mid \mathrm{b}CC \\ C & \longrightarrow & C\mathrm{c} \mid \epsilon \end{array}$$

[15 marks]



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6.

- a. Define *pushdown automaton* (PDA), and explain what it means for a PDA to accept an input. [5 marks]
- b. Describe a *deterministic* pushdown automaton that accepts the set of all palindromes over the alphabet {a, b, c} in which the only place the letter c occurs is at the centre of the word. (For example, words like abcba, aabacabaa.) [10 marks]
- c. Explain why unrestricted palindromes cannot be recognised by deterministic pushdown automata (and one would have to use a nondeterministic PDA instead). [5 marks]

7.

- a. Classify each of the following languages according to whether it is recursively enumerable, recursive, context-free or regular.
 - (i) words over the one-letter alphabet a whose length is a square number [2 marks]
 - (ii) words over the one-letter alphabet a whose length is either a multiple of 3 or a multiple of 4 [2 marks]
 - (iii) encodings of Turing machines that accept the empty string (An encoding of a Turing machine can be assumed to be any standard encoding that lists the states and transitions of the machine.) [2 marks]
- b. Give a detailed description of a Turing machine M that performs a cyclic left shift of its input, assumed to be a string over the alphabet $\{0, 1\}$. So, if the input is of length n, the machine should copy the first letter of the input to position n and copy each of the remaining letters one position to the left. M should halt after performing this operation. For an input of length n, how many transitions would be executed by your machine M, as a function of n? [14 marks]