

## MATHEMATICAL TRIPOS Part III

Tuesday 1 June, 2004 1.30 to 4.30

## PAPER 11

## PROBABILISTIC COMBINATORICS

Attempt FOUR questions.

There are six questions in total. The questions carry equal weight.

You may not start to read the questions printed on the subsequent pages until instructed to do so by the Invigilator. 1 Let  $A_i$ ,  $1 \le i \le n$  be events and let  $J_i \subset [n]$  be such that  $A_i$  is independent of the system  $\{A_j : j \ne i, j \notin J_i\}$ . Let  $p_i = \Pr A_i$  and let  $\Delta = \frac{1}{2} \sum_i \sum_{j \in J_i} \Pr(A \cap B)$ , as usual. Let W be the number of  $A_i$  that occur and let  $\lambda = \mathbb{E}W = \sum_i p_i$ . Prove the Stein-Chen bound:

$$d_{TV}(\mathcal{L}(W), \operatorname{Po}(\lambda)) \le \min(1, \lambda^{-1}) \left[ \sum_{i}^{n} p_{i}^{2} + \sum_{i} p_{i} \sum_{j \in J_{i}} p_{j} + 2\Delta \right]$$

[You may assume without proof the standard inequality  $\Delta g \leq \min(1, \lambda^{-1})$ .]

A random 3-uniform hypergraph of order n is constructed by choosing the edges independently and at random with probability p from the  $\binom{n}{3}$  possible edges. Let W be the number of complete subgraphs of order 4 and let  $\lambda = \mathbb{E}W = \binom{n}{4}p^4$ . Show that

$$d_{TV}(\mathcal{L}(W), \operatorname{Po}(\lambda)) \leq 4np^4 + 4np^3.$$

**2** State and prove the Local Lemma.

Deduce the Ramsey bound  $R(k) \ge (\sqrt{2}/e + o(i)) k 2^{k/2}$ .

The vertices of a cycle of length 12n are partitioned into n sets each of size 12. Show that it is possible to select an independent set of n vertices, one from each set.

**3** Let  $\Omega = \prod_{i=1}^{n} A_i$  be a product probability space and let  $f : \Omega \to \mathbb{R}$  be such that  $|f(\omega) - f(\omega')| \leq c_i$  whenever  $\omega$  and  $\omega'$  differ in only the *i*'th co-ordinate. Prove that  $\Pr\{|f - \mathbb{E}f| \geq t\} \leq 2\exp(-2t^2 / \sum_i c_i^2).$ 

Let  $A \subset \mathcal{P}[n]$  and let  $A_t = \{y \in [n] : \exists x \in A, d(x, y) \leq t\}$  where d(x, y) is the Hamming distance. Prove that, if  $|A| \geq \epsilon 2^n$  and  $t = \sqrt{2n \log(2/\epsilon)}$ , then  $|A_t| \geq (1 - \epsilon)2^n$ .

4 Describe how the semi-random method can be used to prove that, for all  $r \in \mathbb{N}$ and  $\epsilon > 0$ , there exists  $\delta > 0$  so that, for any *r*-uniform hypergraph *G* whose degrees lie between  $(1 - \epsilon)D$  and  $(1 + \epsilon)D$  and whose co-degrees are bounded by  $\delta D$ , the chromatic index  $\chi'(G)$  of *G* satisfies  $\chi'(G) \leq (1 + \epsilon)D$ .

Your essay should make clear the main steps of the proof, and should show how the standard probabilistic tools are applied.

**5** Let  $\Omega = \prod_{i=1}^{n} X_i$  be a product of finite probability spaces. Define the Talagrand distance function  $d_{\mathrm{T}}(x, A)$ , where  $x \in \Omega$  and  $A \subset \Omega$ .

Prove Talagrand's inequality  $\mathbb{E} \exp\{d_{\mathrm{T}}^2(x, A)/4\} \leq 1/\Pr A$ .

Explain how it can be used to prove that the length of a longest increasing subsequence of a random sequence is concentrated near its mean value.

**6** Let  $X = (X_1, \ldots, X_n)$  be a sequence of random variables and, for  $A \subset [n]$ , let  $X_A = (X_i : i \in A)$ . Let  $\mathcal{A} \subset \mathcal{P}[n]$  be such that  $|\{A \in \mathcal{A} : j \in A\}| \geq k$  for all  $j \in [n]$ . Prove that  $kH(X) \leq \sum_{A \in \mathcal{A}} H(X_A)$ , where H is the entropy function. [Standard facts about entropy should be stated if used, but need not be proved.]

A bipartite graph is (a, b)-regular if the vertices in one class have degree a and those in the other class have degree b. Show that an (a, b)-regular bipartite graph of order n has at most  $(2^a + 2^b - 1)^{n/(a+b)}$  independent subsets.