GCE 2005 January Series



# Mark Scheme

# **Computing Specification**

CPT4 Processing and Programming Techniques

Mark schemes are prepared by the Principal Examiner and considered, together with the relevant questions, by a panel of subject teachers. This mark scheme includes any amendments made at the standardisation meeting attended by all examiners and is the scheme which was used by them in this examination. The standardisation meeting ensures that the mark scheme covers the candidates' responses to questions and that every examiner understands and applies it in the same correct way. As preparation for the standardisation meeting each examiner analyses a number of candidates' scripts: alternative answers not already covered by the mark scheme are discussed at the meeting and legislated for. If, after this meeting, examiners encounter unusual answers which have not been discussed at the meeting they are required to refer these to the Principal Examiner.

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### **Computing: Unit CPT4**

#### The following notation is used in the mark scheme

- ; means a single mark;
- / means alternative response;
- A means acceptable creditworthy answer;
- R means reject answer as not creditworthy;
- I means ignore.

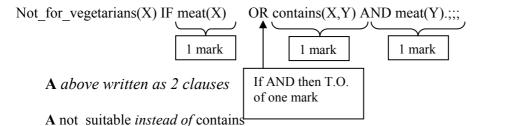
(ii)

I punctuation i.e. accept without . (fullstop) Penalise once in each of (a), (b), (c), (d) for wrong case (must be lower case for predicates and atoms, and variables must be capitalised)

- **1.** (a) (i) contains(raitha, yoghurt).;
  - (b) nuts;eggs; I any punctuation or 'and' I order

dairy product(yoghurt).;

- (c) not\_suitable(X,eggs).;; one mark for not\_suitable(X, ), one mark for 'eggs' in correct place
   A contains (X,eggs);;marking as above
   A not\_suitable(X,eggs) IF contains (X, eggs); give one mark only, as not a goal but a rule
   T.O. if anything else added to correct answer I ?
- (d)



I case for AND. AND can be represented by a comma

**R** contains (X,meat) **R** contains (Meat,X)

accept any variable for Y but must be the same variable in both cases. If not do not give mark for meat(Y).

2.	(a)	system resources exist in separate nodes of a network (with transparent access by users); Tasks can be shared between different nodes; <i>must imply shared processing</i>	у	
		<b>R</b> idea of mainframe and terminals		1
	(b)	network with File server / print server / webserver; A client-server; A an application or specific example e.g. search engines, Internet R brand names		1
	(c)	able to share resources; <b>A</b> <i>data, hardware for resources</i> share processing of task among several processors; local processing done locally; maximise use of resources // faster execution of complex calculation		2
		Tota	al	4

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**3.** (a) (i) *Multi-programming:* concurrent/<u>apparent</u> simultaneous; execution of two (or more) programs; A two (or more) programs seem to be running at the same time;; A job, process, thread, sequence of instructions instead of program

(ii) *Process:* a program currently executing//waiting to be executed an instance of a program: // a program in a phase of execution;
 **R** task/application instead of program

Position	Name	Running Time	Address	Pointer	
1	Process6	7	01400	4 (02300)	1 mark for 4,5,3 correct
3	Process7	17	01700	5 (04100)	1 mark for null
4	Process2	17	02300	3 (01700);	pointer correct
5	Process9	45	04100	-1; A 0;	A sensible const. Name
6	Process5	2	01200	8 (01900)	representing null pointer
8	Process19	5	01900	1 (01400);	<ul> <li>A mark for 8,1 correct</li> </ul>

(b) (i) allow addresses in the Pointer column.

- (ii) array; of records; OR linked list; of records; OR 4 1-D arrays; one for each column; OR one 1-D array for process name; one 2-D arrays for numerical data;2
- (iii) Marks to be allocated as follows:

<i>l for initialisation l for while not at end of list l for printing l for getting <u>next</u> pointer</i>	ListPointer ← <u>HeadPointer;</u> While ListPointer <>-1 Do; Print ListArray[ListPointer]. <u>Name;</u> ListPointer ← ListArray[ListPointer]. <u>Pointer;</u>	Any name acceptable for ListPointer and ListArray
P1 if headpointer is reassigned		
Note: a sorting method gets a maximum	4	

*Alternative solution:* REPEAT UNTIL next=-1 *OR* 

IF listpointer <>-1 then REPEAT....

(iv)

List	Reason
List of suspended/blocked/halted/	waiting for a resource or complete a requested
unrunnable processes;	I/O transfer;
List of inactive/dormant jobs;	Waiting to be admitted to the system;

I currently running processes I interrupt

2

interrupt: a signal; (from a device) seeking the attention of the processor; 2 4. (a) (b) interrupting device supplies; an offset; A index / indexed address R index register added to the base address; A base register *instead of* base address added to the base address; A base register instead of base address gives (start) address of interrupt service routine / ISR / interrupt handler; A location as BOD OR The number supplied by interrupting device; Is an absolute address of a location; That is the (start) address of the ISR / interrupt service routing / interrupt handler;

I anything about priorities A location as BOD

Note: Question is not about how an interrupt is being serviced

5. (a) produces re-useable code because of inheritance/encapsulation; Produces re-useable objects; data is protected // only accessible in well-defined ways (because of encapsulation); more efficient to write programs which use pre-defined / inherited objects / classes;
storage structure of data and method code of a class may be altered without affecting programs that make use of the class; code produced contains fewer errors / more reliable; solutions are easier to understand (when expressed in terms of objects);

easier to enforce design consistency; easier to debug;

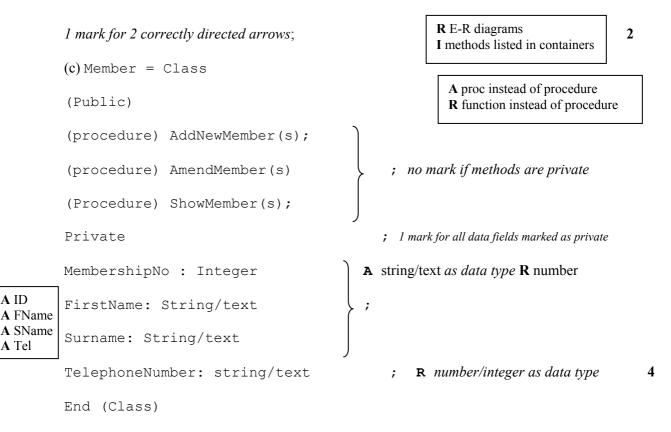
less maintenance effort required by developer since objects can be re-used;

new functions can be added to objects easily (because of inheritance);

**R** Easier to program **I** references to GUIs

2

(b) *I mark for correct base class and derived classes incl. containers;* 



Public may come after Private. Each line may be preceded by Public or Private & in no particular order R diagrammatic answer I case I white space

6.	<ul> <li>(a) <u>the</u> set / list of bit patterns / binary codes representing machine operations; the set / list of bit patterns / binary codes for which machine operations have been defined;</li> <li>The collection of different operations available; A complete set <i>instead of</i> t set</li> <li><i>Note: must imply all available opcodes</i></li> <li><b>R</b> anything that sounds like a program</li> </ul>						
	(b)		a storage unit where <u>datum/data item/instruction</u> is temporarily stored;				
			storage unit that can be <u>accessed rapidly</u> ;	Must imply individual items, eg during FE cycle			
			a special <u>high-speed</u> memory location;	<b>R</b> piece of memory instead of storage unit			
			storage unit internal to the processor;	Answer must distinguish register from buffer			
			storage unit that can be symbolically identified;	1			
	(c)	c) (i) 75 is the number to be copied/stored in the accumulator/load accumulator with value 75;					
		(ii)	The contents of the accumulator is to be copied/stored at address 75;A location <i>inst. of</i> address				
		(iii)	(The contents of location) 75+;(X)/contents of X; ( accumulator)	is to be loaded into the	2		
		(iv)	The contents of the accumulator is to be stored in t address is stored; at address 75;	he location whose			
			A diagrammatic answer		2		
	(d)	(i)	195;		1		
		(ii)	as a shorthand / because it is easier/quicker to read /less likely to make mistakes; Saving space <b>T.O.</b> , saving processor time <b>T.O</b> .	(than lots of 0's and 1's)	1		

(iii)

Program Counter	Content of Index Register X	Content of Accumulator A
00A3	-	-
00A4	1;	-
00A5	1	C3 / 195 / 11000011;
00A6	1∫;	C8 / 200 / 11001000;

Accept trace table moved up 1 row

4

(iv) contents of accumulator / C8 / 200 / 11001000 is stored at address 00A2; 1

7.

(a)		BE4; <i>must be capital letter</i>	S	1
(b)		190.25 / 190 <sup>1</sup> / <sub>4</sub> ;;	one mark for correct integer part,	
			one mark for correct fractional part	
			one mark for correct working	
			(e.g. correct place values)	3
(c)		-1052;;	1 mark for workings if result incorrect	
		1 mark for sign, 1 mark for 105	52	2
(d)	(i)	-	l marks for workings if result incorrect ving binary point 4 places or showing 2 <sup>4</sup>	3
	(ii)	starts with 1 0		
		the first 2 binary digits are diff	erent;	
		a significant bit is stored after the (implied) binary point;		
		bit after (implied) binary point	different from bit before binary point;	1
		A all leading 1's have been rem	noved // there are no leading 1's;	
		<b>R</b> there are no leading zeros	Total	10

## END OF CPT4 MARK SCHEME