

GCE Examinations
Advanced Subsidiary / Advanced Level
Decision Mathematics
Module D2

Paper D

MARKING GUIDE

This guide is intended to be as helpful as possible to teachers by providing concise solutions and indicating how marks should be awarded. There are obviously alternative methods that would also gain full marks.

Method marks (M) are awarded for knowing and using a method.

Accuracy marks (A) can only be awarded when a correct method has been used.

(B) marks are independent of method marks.



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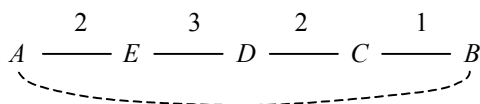
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D2 Paper D – Marking Guide

1. (a)

order:	1	5	4	3	2
	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>
<i>A</i>	–	4	7	8	2
<i>B</i>	4	–	1	5	6
<i>C</i>	7	1	–	2	7
<i>D</i>	8	5	2	–	3
<i>E</i>	2	6	7	3	–

M1



A1

upper bound = $2 \times$ weight of MST
 $= 2 \times (2 + 3 + 2 + 1) = 2 \times 8 = 16$ miles

M1 A1

- (b) use *AB* saving $2 + 3 + 2 + 1 - 4 = 4$
 new upper bound = $16 - 4 = 12$ miles

M1

A1

(6)

2. (a) adding 5 to all entries to make them positive gives

M1

		<i>B</i>		
		I	II	III
<i>A</i>	I	11	1	4
	II	3	10	8
	III	10	6	2

new value of game $v = V + 5$

A1

- (b) let *B* play strategies I, II and III with proportions p_1, p_2 and p_3
 let $x_1 = \frac{p_1}{v}, x_2 = \frac{p_2}{v}, x_3 = \frac{p_3}{v}$

M1

A1

- (c) $p_1 + p_2 + p_3 = 1$
 dividing by v gives $x_1 + x_2 + x_3 = \frac{1}{v}$
 we wish to minimise $v \therefore$ maximise $\frac{1}{v}$
 objective function is maximise $P = x_1 + x_2 + x_3$

M1

A1

- (d) from *A* I, $11p_1 + p_2 + 4p_3 \leq v$
 from *A* II, $3p_1 + 10p_2 + 8p_3 \leq v$
 from *A* III, $10p_1 + 6p_2 + 2p_3 \leq v$

M1

dividing by v gives the constraints

$$\begin{aligned}
 11x_1 + x_2 + 4x_3 &\leq 1 \\
 3x_1 + 10x_2 + 8x_3 &\leq 1 \\
 10x_1 + 6x_2 + 2x_3 &\leq 1 \\
 \text{also } x_1 \geq 0, x_2 \geq 0, x_3 \geq 0
 \end{aligned}$$

A1

(8)

3. (a)

order:	1	6	2	3	5	4	7
	A	B	C	D	E	F	G
A	–	83	57	68	103	91	120
B	83	–	78	63	41	82	52
C	57	78	–	37	59	63	74
D	68	63	37	–	60	52	62
E	103	41	59	60	–	48	51
F	91	82	63	52	48	–	77
G	120	52	74	62	51	77	–

M1 A1

tour: *ACDFEBGA*

upper bound = $57 + 37 + 52 + 48 + 41 + 52 + 120 = 407$ miles

A1

A1

(b) e.g. starting at *B*

order:	1	6	5	2	3	4	
	A	B	C	D	E	F	G
A	–	83	57	68	103	91	120
B	83	–	78	63	41	82	52
C	57	78	–	37	59	63	74
D	68	63	37	–	60	52	62
E	103	41	59	60	–	48	51
F	91	82	63	52	48	–	77
G	120	52	74	62	51	77	–

M1 A1

lower bound = weight of MST + two edges of least weight from *A*

= $(41 + 48 + 51 + 52 + 37) + 57 + 68 = 354$ miles

M1 A1

(c) $354 \leq d \leq 407$

B1

(9)

4.

Stage	State	Action	Destination	Value
1	<i>I</i>	<i>IL</i>	<i>L</i>	5*
	<i>J</i>	<i>JL</i>	<i>L</i>	6*
	<i>K</i>	<i>KL</i>	<i>L</i>	10*
2	<i>F</i>	<i>FI</i>	<i>I</i>	$\min(5, 5) = 5^*$
		<i>FJ</i>	<i>J</i>	$\min(2, 6) = 2$
		<i>FK</i>	<i>K</i>	$\min(2, 10) = 2$
	<i>G</i>	<i>GI</i>	<i>I</i>	$\min(8, 5) = 5$
		<i>GJ</i>	<i>J</i>	$\min(9, 6) = 6^*$
		<i>GK</i>	<i>K</i>	$\min(3, 10) = 3$
<i>H</i>	<i>HI</i>	<i>I</i>	$\min(10, 5) = 5$	
	<i>HJ</i>	<i>J</i>	$\min(2, 6) = 2$	
	<i>HK</i>	<i>K</i>	$\min(9, 10) = 9^*$	
3	<i>B</i>	<i>BF</i>	<i>F</i>	$\min(8, 5) = 5$
		<i>BG</i>	<i>G</i>	$\min(11, 6) = 6^*$
		<i>BH</i>	<i>H</i>	$\min(4, 9) = 4$
	<i>C</i>	<i>CF</i>	<i>F</i>	$\min(5, 5) = 5$
		<i>CH</i>	<i>H</i>	$\min(10.5, 9) = 9^*$
	<i>D</i>	<i>DF</i>	<i>F</i>	$\min(9, 5) = 5$
		<i>DH</i>	<i>H</i>	$\min(6, 9) = 6^*$
	<i>E</i>	<i>EF</i>	<i>F</i>	$\min(12, 5) = 5$
<i>EG</i>		<i>G</i>	$\min(7, 6) = 6$	
<i>EH</i>		<i>H</i>	$\min(15, 9) = 9^*$	
4	<i>A</i>	<i>AB</i>	<i>B</i>	$\min(1, 6) = 1$
		<i>AC</i>	<i>C</i>	$\min(4.5, 9) = 4.5$
		<i>AD</i>	<i>D</i>	$\min(13, 6) = 6$
		<i>AE</i>	<i>E</i>	$\min(10, 9) = 9^*$

A1

M1 A2

M1 A1

A1

giving route *AEHKL*
shortest stage is 9 miles

M1 A1

A1

(10)

5. need to add dummy column giving

M1

19 69 168 0
22 64 157 0
20 72 166 0
23 66 171 0

col min. 19 64 157 0

reducing rows will make no difference

B1

reducing columns gives:

~~0~~ ~~5~~ ~~11~~ ~~0~~
~~3~~ ~~0~~ ~~0~~ ~~0~~
1 8 9 0
4 2 14 0

(N.B. a different choice of lines will lead to the same final assignment)

M1 A1

3 lines required to cover all zeros, apply algorithm

B1

~~0~~ ~~5~~ ~~11~~ ~~1~~
~~3~~ ~~0~~ ~~0~~ ~~1~~
~~0~~ ~~7~~ ~~8~~ ~~0~~
~~3~~ ~~1~~ ~~13~~ ~~0~~

M1 A1

3 lines required to cover all zeros, apply algorithm

~~0*~~ ~~4~~ ~~10~~ ~~1~~
~~4~~ ~~0~~ ~~0*~~ ~~2~~
~~0~~ ~~6~~ ~~7~~ ~~0*~~
~~3~~ ~~0*~~ ~~12~~ ~~0~~

A1

4 lines required to cover all zeros so allocation is possible

B1

stage 1 is run by Alex

stage 2 is run by Suraj

stage 3 is run by Darren

Leroy does not take part

M1 A1 (11)

6. (a)

		Y		row minimum
		Y_1	Y_2	
X	X_1	-2	4	-2
	X_2	6	-1	-1
column maximum		6	4	

M1 A1

$\max(\text{row min}) = -1$ $\min(\text{col max}) = 4$
 $\max(\text{row min}) \neq \min(\text{col max}) \therefore$ no saddle point

B1

(b) (i) let X play strategies X_1 and X_2 with proportions p and $(1 - p)$
 expected payoff to X against each of Y 's strategies:

$$\begin{aligned}
 Y_1 \quad & -2p + 6(1 - p) = 6 - 8p \\
 Y_2 \quad & 4p - (1 - p) = 5p - 1
 \end{aligned}$$

M1 A1

$$\begin{aligned}
 \text{for optimal strategy} \quad & 6 - 8p = 5p - 1 \\
 & \therefore 13p = 7, \quad p = \frac{7}{13}
 \end{aligned}$$

$\therefore X$ should play X_1 $\frac{7}{13}$ of time and X_2 $\frac{6}{13}$ of time

M1 A1

(ii) let Y play strategies Y_1 and Y_2 with proportions q and $(1 - q)$
 expected loss to Y against each of X 's strategies:

$$\begin{aligned}
 X_1 \quad & -2q + 4(1 - q) = 4 - 6q \\
 X_2 \quad & 6q - (1 - q) = 7q - 1
 \end{aligned}$$

M1 A1

$$\begin{aligned}
 \text{for optimal strategy} \quad & 4 - 6q = 7q - 1 \\
 & \therefore 13q = 5, \quad q = \frac{5}{13}
 \end{aligned}$$

$\therefore Y$ should play Y_1 $\frac{5}{13}$ of time and Y_2 $\frac{8}{13}$ of time

M1 A1

(c) value of game = $6 - (8 \times \frac{7}{13}) = 1 \frac{9}{13}$

M1 A1 (13)

7. (a)

	D	E	F	Available
A	20			20
B	10	5		15
C			25	25
Required	30	5	25	

M1 A1

no. of rows + no. of cols - 1 = 3 + 3 - 1 = 5

in this solution only 4 cells are occupied, less than 5 ∴ degenerate

B1

(b) placing 0 in (3, 2) as it has lowest cost of unoccupied cells

taking $R_1 = 0$, $R_1 + K_1 = 13$ ∴ $K_1 = 13$ $R_2 + K_1 = 10$ ∴ $R_2 = -3$

$R_2 + K_2 = 9$ ∴ $K_2 = 12$ $R_3 + K_2 = 6$ ∴ $R_3 = -6$

$R_3 + K_3 = 8$ ∴ $K_3 = 14$

M1 A2

	$K_1 = 13$	$K_2 = 12$	$K_3 = 14$
$R_1 = 0$	0	11	14
$R_2 = -3$	0	0	12
$R_3 = -6$	15	0	0

improvement indices, $I_{ij} = C_{ij} - R_i - K_j$

∴ $I_{12} = 11 - 0 - 12 = -1$

$I_{13} = 14 - 0 - 14 = 0$

$I_{23} = 12 - (-3) - 14 = 1$

$I_{31} = 15 - (-6) - 13 = 8$

M1 A1

pattern not optimal as there is a negative improvement index

B1

applying algorithm

let $\theta = 5$, giving

	D	E	F
A	$20 - \theta$	θ	
B	$10 + \theta$	$5 - \theta$	
C			25

	D	E	F
A	15	5	
B	15		
C			25

M1 A1

this solution is also degenerate

place 0 in (3, 2) again

taking $R_1 = 0$, $R_1 + K_1 = 13$ ∴ $K_1 = 13$ $R_1 + K_2 = 11$ ∴ $K_2 = 11$

$R_2 + K_1 = 10$ ∴ $R_2 = -3$ $R_3 + K_2 = 6$ ∴ $R_3 = -5$

$R_3 + K_3 = 8$ ∴ $K_3 = 13$

M1 A1

	$K_1 = 13$	$K_2 = 11$	$K_3 = 13$
$R_1 = 0$	0	0	14
$R_2 = -3$	0	9	12
$R_3 = -5$	15	0	0

∴ $I_{13} = 14 - 0 - 13 = 1$

$I_{22} = 9 - (-3) - 11 = 1$

$I_{23} = 12 - (-3) - 13 = 2$

$I_{31} = 15 - (-5) - 13 = 7$

M1 A1

all improvement indices are non-negative ∴ pattern is optimal

B1

15 units from A to D, 5 units from A to E,

15 units from B to D, 25 units from C to F

A1

total cost = $(15 \times 13) + (5 \times 11) + (15 \times 10) + (25 \times 8) = \text{£}600$

A1

(18)

Total

(75)

Performance Record – D2 Paper D

Question no.	1	2	3	4	5	6	7	Total
Topic(s)	TSP, shortcuts	game, formulate lin. prog.	TSP, nearest neighbour	dynamic prog., maximin	allocation, dummy	game	transport., n-w corner, stepping-stone, degeneracy	
Marks	6	8	9	10	11	13	18	75
Student								